



Reactive Power Optimization in Grid using Imperialist Competitive Algorithm

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ABSTRACT: The Imperialist Competitive Algorithm (ICA) is a heuristic technique for optimization procedure. This algorithm is based on socio-politically inspired optimization strategy. It has shown better convergence rate to achieve global optimum for different optimization problems. The IEEE 30-bus test system is used to demonstrate some aspects of the application of the ICA on the reactive power market problem. The outcomes are compared with those obtained by the Genetic Algorithm (GA) showing that the ICA convergence rate is more than GA.

KEYWORDS: Imperialist Competitive Algorithm, reactive power optimization.

I. INTRODUCTION

Reactive power optimization is a short-term planning activity carried out by system operators in order to ensure secure power system operation. Reactive power has a significant effect on the system security as it is directly associated with the power system voltage stability [1]. The reactive power can be controlled in order to improve the voltage profile and to minimize the system loss. Generally some load bus voltage may violate their upper or lower limits during system operation due to the disturbance or system configuration changes. The power system operator can alleviate this situation and voltages can be maintained within their permissible limits by reallocating reactive power generation in the system [2]. This means by adjusting the generator voltages, transformer taps and switchable VAR sources (capacitor or reactor).

In this paper, Imperialist Competitive Algorithm (ICA) is a population based algorithm, used to solve the reactive power optimization problem in a power grid. Conjointly, implementation of the ICA on reactive power optimization problem is compared with Genetic Algorithm.

This paper put forwards the efficient communication between CR nodes and spectrum utilization. Secondly, the security concerns of spectrum sensing are to ensure steadfastness. It uses two selection schemes called Node Selection Scheme (NSS) and Channel Selection Scheme (CSS). The aim of NSS is to allow each node to check its gain in copying a message to a relay while reconnoitring its transmission effort. Using NSS, each node decides which path should be used in order to provide minimum energy consumption without enduring end-to-end delay performance. Based on CSS, each node decides and switches to a licensed channel to maximize spectrum utilization while keeping the interference in a minimum level. This eventually enables CR-Network nodes to determine optimum path nodes and channels for an efficient communication in CR-Networks. The CR technology allows Secondary Users (SUs) to seek and utilize “spectrum holes” in a time and location-varying radio environment without causing harmful interference to Primary Users (PUs). This opportunistic use of the spectrum, leads to new challenges to the varying available spectrum. Using a Trust-Worthy algorithm, it improves the firmness of the Spectrum sensing in CR-Networks.

II. IMPERIALISTIC COMPETITIVE ALGORITHM

The Imperialist Competitive Algorithm (ICA) that has been recently introduced is used here to solve the reactive power optimization problem. This algorithm is based on real-world countries trying to extend their power over other countries in order to use their resources. In other words, imperialist countries attempt to dominate other countries and compete strongly with each other for taking possession of other countries. During this competition, robust empires become

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powerful and the weakest one will collapse. The ICA starts with an initial population called countries. Some of the best countries in the population are selected to be the imperialists and the rest form the colonies of these imperialists. Each imperialist with his colonies forms the empire. The more powerful imperialist, have the more colonies. In other words, the imperialist with a larger empire is more powerful. When the imperialist competition starts, their attempt to achieve more colonies starts. And the colonies start to move towards their imperialists. At the end, just one imperialist will remain.

The steps used in the ICA are as follows:

1. Form the initial population of decision variables.
2. Select some random points on the function and initialize the empires.
3. Move the colonies toward their relevant imperialist (Assimilation).
4. Randomly change the position of some colonies (Revolution).
5. If there is a colony in an empire which has lower cost than the imperialist, exchange the positions of that colony and the imperialist.
6. Unite the similar empires.
7. Compute the total cost of all empires.
8. Pick the weakest colony (colonies) from the weakest empires and give it (them) to one of the empires (Imperialistic competition).
9. Eliminate the powerless empires.
10. It terminates if stop condition convinced, else go to step 2.

More details about the ICA can be obtained from [9]. The Imperialistic Competitive Algorithm is a newly introduced global search strategy that has recently been implemented for solving different optimization problems. This evolutionary optimization technique has shown great performance in both convergence rate and global optimum achievement. The initial population consists of individuals named country in the beginning of the algorithm and then the populations are separated to two different types as imperialist and colony based on their power (fitness value). Some of the best initial countries (the countries with the least fitness value) become Imperialists and start taking control over other countries (called Colonies) and form the initial Empires. The stage of this algorithm is conferred below.

Initial empires formation

At the start of the algorithm, an initial population (countries) should be created. Each country consists of N variables to be optimized. The Fig. 1 shows the characteristic of empires formation.

$$\text{Country}_i = [x_{1,i} \quad x_{2,i} \quad x_{3,i} \quad \dots \quad x_{Nvar,i}]$$

The cost function of countries is calculated by evaluating the function f at the variables.

$$\text{Cost}_i = f(\text{Country}_i) = f(x_{1,i} \quad x_{2,i} \quad x_{3,i} \quad \dots \quad x_{Nvar,i})$$

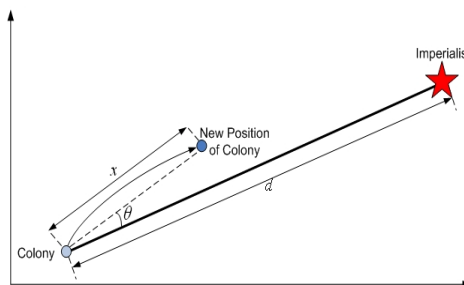


Fig.1 Empires formation

In the beginning, $N_{country}$ is defined as the initial number of empire countries. The best countries of this empire are selected as imperialists (N_{imp}). Each of the remaining countries is under the control of one of those imperialists as colonies (N_{col}). These colonies are divided between imperialists depending upon their power. In order to perform this work, first we should calculate the normalized cost of imperialists using equation below

$$C_n = \max \{c_i\} - c_n$$

Where, C_n is the cost of n th imperialist, $\max \{C_i\}$ is maximum cost of imperialists and $n C$ is the normalized cost of n th imperialist. Using normalized costs, the normalized power of imperialists can be calculated.

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$$P_n = \left| \frac{C_n}{\sum_{i=1}^{N_{imp}} C_i} \right|$$

Finally the initial number of colonies that belong to each imperialist is calculated:

$$N.C._n = \text{round} \{ p_n \cdot (N_{col}) \}$$

Where, $N.C._n$ is the initial number of colonies that belong to an imperialist, N_{col} is the number of all colonies in initial population and round is a function which returns the nearest integer to a floating point number.

Movement of colonies towards the imperialists

In this stage of algorithm, the colonies of each imperialist are moved towards them. If the problem has n variables, there would be an n dimensional space in which each country is a point on this space. The colony is moved by x units towards the imperialist. 'x' is a random number in the interval $(0, \beta d)$ which usually has uniform distribution. 'd' is considered as the distance between the imperialist and the colony and β is a number larger than one and nearly two. $\beta=2$ is a good selection. Movement is not mandatory in the direction of the line towards the imperialist and it can deviate from this line (Figure 1). The deviation (θ) depends upon the problem and should be tuned. $\theta=\pi/4$ is a optimum selection for many applications. Equation shows this process in two dimensions:

$$\text{Newcol}_n^i = \text{col}_n^i + \beta * \text{rand} * (\text{imp}_n - \text{col}_n^i)$$

Where, col_n^i is the position of the i^{th} colony of the n^{th} imperialist, rand is a random number in $(0, 1)$ and imp_n is the position of the n^{th} imperialist. After the movement of the colonies towards the imperialist, if the colonies reach a better position than imperialist (with a lower fitness value) their position would be switched. The Fig. 2 shows the exchanging the positions of a colony and Imperialist

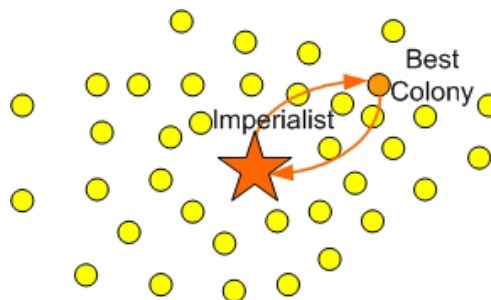


Fig. 2 Exchanging the positions of a colony and Imperialist

Imperialistic Competition and elimination of weak empires

In this stage, the weakest colony of the weakest imperialist is given to other imperialists depending on their power (probability). For calculating this probability, first we should calculate the total cost of the empires. The total cost of an empire equals to the sum of its own cost and a part of cost of its colonies:

$$T.C._n = \text{Cost}(\text{imperialist}) + \zeta \text{ mean} \{ \text{Cost}(\text{colonies of empires}) \}$$

Where, $T.C._n$ is the total cost of n^{th} empire and ζ is a positive number which is considered to be less than 1. This number represents the effect of colonies in the total cost of empire. Then we can calculate the normalized total cost of empire:

$$N.T.C._n = \max \{ T.C._i \} - T.C._n$$

Where, $N.T.C._n$ is the normalized total cost of n^{th} empire. Finally we can calculate the probability (power) of each empire for absorbing the weakest colony.

$$P_{pn} = \left| \frac{N.T.C._n}{\sum_{i=1}^{N_{imp}} N.T.C._i} \right|$$

More powerful imperialists have more chance to absorb the weakest colony. For implementation of this issue, the following works should be done. First we should form P_p and R matrices:

$$P_p = [P_{p1}, P_{p2}, \dots, P_{pN_{imp}}]$$

$$R = [r_1, r_2, \dots, r_{N_{imp}}]$$

Where, P_p is the matrix of probability and r_i is a random number in $(0, 1)$. Then D vector is formed.

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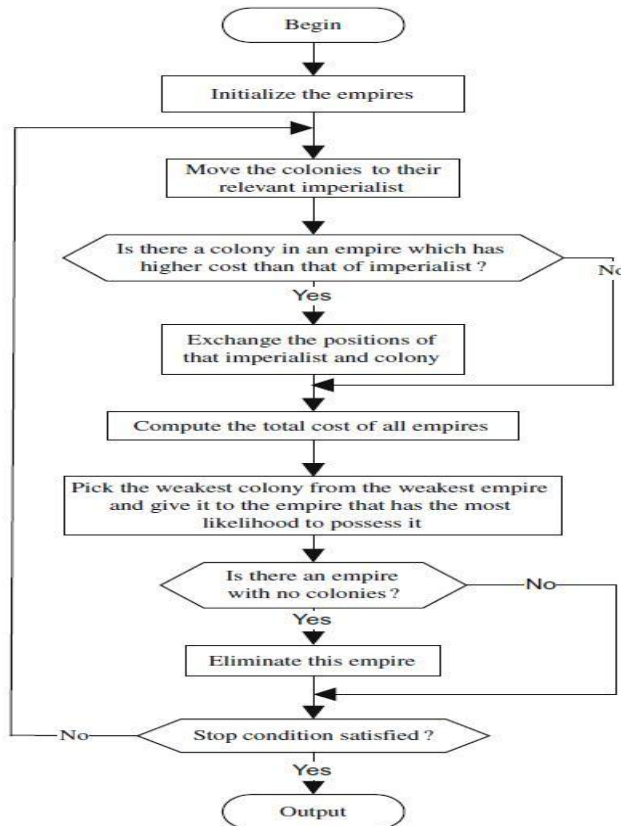
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$$D = P_p - R = [P_{p1-r_1}, P_{p2-r_2}, \dots, P_{pN_{imp}-r_{N_{imp}}}]$$

The imperialist which has the maximum relevant index in D vector will absorb the mentioned colony. If an imperialist loses all of its colonies, it will collapse.

Imperial flow representation



Convergence

After a while, all empires will collapse except the most powerful one and all the other countries will be under the control of this empire. At this point, the algorithm stops. This condition cannot be ample enough for obtaining accurate results. Sometimes, it is required that the position of some or all the countries would be the same which is an ideal solution. And sometimes the maximum iteration number can be considered as a criterion for stopping the algorithm. The weakest colony in empire is shown in Fig.3. In this method weakest colony in empire is easily found.

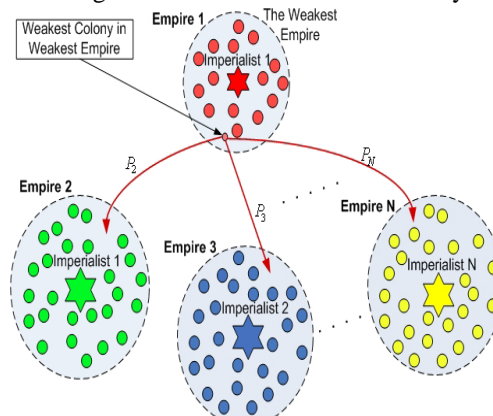


Fig. 3 Weakest colony in empire



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III. PROBLEM FORMULATION

The problem of reactive power optimization is directly concerned not only with the service quality and reliability of supply but also with the economy and security of the power systems. Generally, some load bus voltage might violate their upper or lower limits during system operation due to disturbance and system configuration changes [3].

Objective Function

The primary objective of the Reactive Power Optimization (RPO) problem is to reduce the system's real power losses and to obtain the setting of various controls for the same. The real power loss is a non-linear function of bus voltage magnitudes and phase angles, which are implicitly a function of the control variables [4].

The objective function is

$$\text{Min } f(x) = \sum_{k=1}^{nl} P_{kLoss} = \sum_{k=1}^{nl} G_k [V_i^2 + V_j^2 - 2V_i V_j \cos \theta_{ij}]$$

Where,

- P_{Loss} Real power loss in the transmission network
- G_k Conductance of the line connecting i and j bus
- V_i Voltage magnitude on the bus i
- V_j Voltage magnitude on the bus j
- θ_{ij} Phase angle of the voltage value between i and j bus

Equality Constraints

$$P_{gi} - P_{di} - V_i \sum_{j=1}^{n_b} V_j (G_{ij} \cos \theta_{ij} + B_{ij} \sin \theta_{ij}) = 0$$

$$Q_{gi} - Q_{di} - V_i \sum_{j=1}^{n_b} V_j (G_{ij} \sin \theta_{ij} - B_{ij} \cos \theta_{ij}) = 0$$

Where,

- P_g Active power of the generation
- Q_g Reactive power of the generation
- P_d Demand of the active power
- Q_d Demand of the reactive power
- G_{ij} Conductance of the line ith and jth bus
- B_{ij} Susceptance of the line ith and jth bus
- θ_{ij} Voltage angle between bus i and bus j
- n_b Number of buses

Inequality Constraints

$$Q_{gi} \min < Q_{gi} < Q_{gi} \max_{i \in n_g}$$

$$V_i \min < V_i < V_i \max_{i \in n_b}$$

$$T_k \min < T_k < T_k \max_{k \in n_t}$$

$$Q_{ci} \min < Q_{ci} < Q_{ci} \max_{i \in n_c}$$

$$V_i \geq V_{i \max} \rightarrow V_i = V_{i \max}$$

$$V_i \leq V_{i \min} \rightarrow V_i = V_{i \min}$$

$$Q_i \geq Q_{i \max} \rightarrow Q_i = Q_{i \max}$$

$$Q_i \leq Q_{i \min} \rightarrow Q_i = Q_{i \min}$$

Power flow equations are used as equality constraints, reactive power source installation restrictions, reactive power generation restrictions, transformer tap-setting and bus voltage restrictions are used as inequality constraints [5].

In the reactive power optimization problem, the generator bus voltages, the tap position of transformer, the amount of reactive power source installations [6] are control variables which are self-constrained. Voltages of load buses and injected reactive power of generator buses are constrained by adding them as penalty terms to the objective function.



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IV. OPTIMIZATION RESULTS

	Fuzzy Logic [3]	PSO [8]	Bacterial Foraging [10]	Proposed Method
Total kVAR	2400	2063	2550	2400
Real Power loss (KW)	165.07	168.80	161.86	160.92
Reactive Power Loss (KVA)	47.825	48.961	47.308	47.047
Min Voltage	0.9515	0.9496	0.9500	0.9507

The above table shows the result comparison of fuzzy logic, PSO, Bacterial foraging and proposed method of total kVAR, real power loss, reactive power loss and minimum voltage. It is evident that the proposed method has minimum losses and better voltage profile compared to methods listed above.

V. CONCLUSION

This technique is designed for the distributed grid network to reduce the real power loss by selecting the minimum loss path implementing the ICA at the network. These generated networks are compared with PSO, Fuzzy Logic, Bacterial Foraging and other obtained values with lesser iteration and has lesser time.

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